## NONCREATIVITY AND TRANSLATABILITY IN TERMS OF INTENSION (\*)

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I intend to show in this paper that Carnap's notion of intension (1) may have interesting applications in logical inquiries. The notion of intension will be applied here to the problem of extending of theories. I will only consider the intensions of sentences (without free variables). The intensions of terms and of other formulas will not be taken into account here (2).

The symbols  $\sim$ ,  $\vee$ ,  $\wedge$ ,  $\equiv$  will be used for the sentential connectives (negation, alternative, conjunction and biconditional).



Let  $T_1$  and  $T_2$  be two standard formalized theories (3) such that  $T_1$  is a subtheory of  $T_2$ . It means that the language  $L_1$  of  $T_1$  is a sublanguage of the language  $L_2$  of  $T_2$  and  $T_1$  is a subset of  $T_2$ . We may suppose that there exists a set  $A_1$  of sentences in  $L_1$  and a set  $A_2$  of sentences in  $L_2$  such that the following equivalence holds for each sentence p in  $L_K$  where k = 1, 2:

the sentence p is in  $T_K$  i.e. p is a theorem of  $T_K$  if and only if p follows logically from  $A_K$ .

The assumption that  $T_1$  is a subset of  $T_2$  is equivalent to the following condition:

- (\*) English translation of a lecture given in Russian at the Faculty of Mathematics of the University of Sofia, June 6 1966, Sofia, Bulgaria.
- (1) R. CARNAP, *Meaning and Necessity*, Chicago 1947. See also my paper An Essay in the Formal Theory of Extension and of Intension, *Studia Logica*, vol. 20, in print.
- (2) A more extended paper concerning the applications of the notion of intension will be published elsewhere.
- (3) I.e. the theories T<sub>1</sub> and T<sub>2</sub> are based on the engere Praedikaten-kalkül mit Identität, Funktionszeichen und Beschreibungssymbolen. See D. HILBERT und P. BERNAYS, *Grundlagen der Mathematik*, vol. I, Berlin 1943.

every sentence in  $A_1$  follows logically from  $A_2$ .

The symbols  $T_1/T_2$  means the extending transformation of  $T_1$  to  $T_2$ . The transformation  $T_1/T_2$  may be *noncreative* (St. Lesniewski) and/or *translatable* (K. Adjukiewicz). The noncreativity and translatability are defined as follows (4).

The transformation  $T_1/T_2$  is noncreative if and only if every sentence in  $L_1$  being a theorem of  $T_2$  is also a theorem of  $T_1$ .

The transformation  $T_1/T_2$  is translatable if and only if for every sentence p in  $L_2$  there exists a sentence q in  $L_1$  such that the biconditional  $p \equiv q$  is a theorem of  $T_2$ .

We say that the transformation  $T_1/T_2$  is a definitional one if and only if

- 1)  $L_2 = L_1 + certain new extralogical constants and$
- 2)  $A_2 = A_1 +$ the standard definitions of all new extralogical constants.

We know very well that if the fransformation  $T_1/T_2$  is definitional then it is both non creative and translatable.



Let k = 1,2. If p is a sentence in  $L_K$  then  $\operatorname{int}_K(p) = \operatorname{the}$  intension of p in  $T_K$ . The intensions of sentences in  $L_K$  may be introduced as follows:

 $int_K(p)=int_K(q)$  if and only if the biconditional  $p\equiv q$  is a theorem of  $T_K$ .

We know very well that the set  $I_K$  of all intensions of sentences in  $L_K$  is a Boolean algebra (called also Lindenbaum algebra of  $T_K$  in  $L_K$ ) such that for all p, q in  $L_K$ :

- 1)  $int_K(p \lor \sim p) = the unit-element in I_K$ ,
- 2)  $int_K(p \land \sim p) = the zero-element in I_K$ ,
- 3)  $-int_K(p) = int_K(\sim p)$ ,
- 4)  $int_K(p) + int_K(q) = int_K(p \vee q)$ ,
- 5)  $\operatorname{int}_{K}(p)$ .  $\operatorname{int}_{K}(q) = \operatorname{int}_{K}(p \wedge q)$ .

The signs -, +, . stand for the Boolean operations of complementation, of addition and of multiplication, respectively.

(4) See K. ADJUKIEWICZ, Die Definition, Actes du Congrès International de Philosophie Scientifique (Sorbonne 1935), vol. II, 1-7, Paris, 1936.

The algebras  $I_1$  and  $I_2$  are quite different, in general. Similarly, if p is a sentence in  $L_1$  then  $int_1(p)$  and  $int_2(p)$  are different, in general. However, from the assumption that  $T_1$  is a subtheory of  $T_2$  it follows that there exists a special connection between  $I_1$  and  $I_2$ .

The intensions of sentences in  $L_2$  may be divide into old and new. If p is a sentence in  $L_2$  then  $int_2(p)$  is called old if and only if there exists a sentence q in  $L_1$  such that  $int_2(p) = int_2(q)$ . Otherwise,  $int_2(p)$  is called new.

It is clear that  $int_2(p)$  is old for each sentence p in  $L_1$ . Of course,  $int_2(p)$  may be new for some sentences p in  $L_2$  only.

Let  $I_{2}^{+}$  denote the set of all old intensions of sentences in  $L_{2}$ . Clearly,  $I_{2}^{+}$  is a subset of  $I_{2}$ . Moreover,  $I_{2}^{+}$  is a Boolean algebra and a subalgebra of  $I_{2}$ . It means that for all p, q in  $L_{2}$ :

- 1) int<sub>2</sub>(p  $\vee \sim$  p) and int<sub>2</sub>(p  $\wedge \sim$  p) belong to I+<sub>2</sub> and
- 2) if  $int_2(p)$  and  $int_2(q)$  belong to  $I^+_2$  then  $-int_2(p)$ ,  $int_2(p) + int_2(q)$  and  $int_2(p)$ .  $int_2(q)$  belong also to  $I^+_2$ .

Consider now the following functional correspondance H between the elements of  $I_1$  and those of  $I_2$ :

$$H(int_1(p)) = int_2(p)$$

for every sentence p in  $L_1$ . Clearly,  $I_1$  = the set of all arguments of H and  $I^+_2$  = the set of all values of H.

The function H is a homomorphism of  $I_1$  onto  $I^{+}_2$  and, consequently, a homomorphism of  $I_1$  into  $I_2$ . This means that for all p, q in  $L_1$ :

- 1)  $H(-int_1(p)) = -H(int_1)p)$ ,
- 2)  $H(int_1(p) + int_1(q)) = H(int_1(p)) + H(int_1(q)),$
- 3)  $H(int_1(p) \cdot int_1(q)) = H(int_1(p)) \cdot H(int_1(q))$ .



An one-to-one homomorphism is called isomorphism. Therefore, the function H is an isomorphism between  $I_1$  and  $I_2$  if and only if the following condition holds for all p, q in  $L_1$ :

(\*) if 
$$int_1(p) \neq int_1(q)$$
 then  $int_2(p) \neq int_2(q)$ .

Theorem 1. The extending transformation  $T_1/T_2$  is noncreative if and only if the function H is an isomorphism between  $I_1$  and  $I_2$ .

Proof. (1) Let the transformation  $T_1/T_2$  be noncreative. Suppose that  $int_2(p) = int_2(q)$  where p, q are in  $L_1$ . Consequently, the biconditional  $p \equiv q$  is a theorem of  $T_2$ . It follows from the noncreativity that the biconditional  $p \equiv q$  is also a theorem of  $T_1$ . This means that  $int_1(p) = int_1/q/$ . The condition (\*) holds. (2) Let the transformation  $T_1/T_2$  be creative. I.e. there exists a sentence p in  $L_1$  such that p is a theorem of  $T_2$  but it is not a theorem of  $T_1$ . It follows that the biconditional  $p \equiv (p \lor \sim p)$  is in  $T_2$  but it is not in  $T_1$ . Consequently,  $int_1(p) \neq int_1(p \lor \sim p)$  and  $int_2(p) = int_2(p \lor \sim p)$ . The condition (\*) does not hold.

It is easy to see that the transformation  $T_1/T_2$  is translatable if and only if  $int_2(p)$  is old for every sentence p in  $L_2$ , i.e. if the new intensions of sentences in  $L_2$  do not exist  $(I_2 = I_2)$ . Thus we have

Theorem 2. The extending transformation  $T_1/T_2$  is translatable and noncreative if and only if the function H is an isomorphism between  $I_1$  and  $I_2$ .

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